
What's new in GHC 7.8

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Type holes

A type hole can be placed wherever a Haskell expression should go.

It tells you:

- Which type the compiler is expecting where the hole is;
 - The types of bound variables.
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Type holes

`(.) :: (b -> c) -> (a -> b) -> (a -> c)`

`f . g = _`

Found hole `'_'` with type: `a -> c`

Where: `'c'` is a rigid type variable bound by ...

`'a'` is a rigid type variable bound by ...

Relevant bindings include

`g :: a -> b (bound at ...)`

`f :: b -> c (bound at ...)`

newtype coercions

The problem:

```
newtype Age = Age Int
```

```
toAgeMap :: HashMap Int String -> HashMap Age String
```

```
toAgeMap = ?
```

newtype coercions

```
coerce :: Coercible a b => a -> b
```

```
newtype Age = Age Int
```

```
toAgeMap :: HashMap Int String -> HashMap MyInt String
```

```
toAgeMap = coerce
```

newtype coercions

Coercible is not a normal type class, its rules being hard-coded in GHC:

- $\forall a. \text{Coercible } a \ a$
 - $\forall a \ b. \text{Coercible } a \ b \Rightarrow \text{Coercible } b \ a$
 - $\forall a \ b \ c. (\text{Coercible } a \ b, \text{Coercible } b \ c) \Rightarrow \text{Coercible } a \ c$
 - $\forall a. \text{given newtype } C \ t_1 \ \dots \ t_n = C \ a \Rightarrow \text{Coercible } a \ (C \ t_1 \ \dots \ t_n)$
 - $\forall a_1 \ \dots \ a_n \ b_1 \ \dots \ b_n. (\text{Coercible } a_1 \ b_1, \dots, \text{Coercible } a_n \ b_n) \Rightarrow \text{Coercible } (C \ a_1 \ \dots \ a_n) \ (C \ b_1 \ \dots \ b_n)$
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newtype coercions

```
newtype Ptr a = Ptr Addr#
```

```
Coercible (Ptr CInt) (Ptr CString)?
```

Type families

```
type family Key (f :: * -> *) :: *
class Functor f => Lookup f where
  lookup :: f a -> Key f -> Maybe a
```

```
type instance Key [] = Int
instance Lookup [] where
  lookup xs ix | ix < 0 = Nothing
  lookup [] _           = Nothing
  lookup (x : xs) ix    =
    if ix == 0 then Just x
    else lookup xs (ix - 1)
```

```
type instance Key (Map k) = k
instance Lookup (Map k) where
  lookup = Map.lookup
```

Closed type families

Type families are open, and thus overlapping definitions are forbidden:

```
{-# LANGUAGE DataKinds, TypeFamilies #-}
```

```
data Nat = Zero | Succ Nat
```

```
type family CountArgs (f :: *) :: Nat
```

```
type instance CountArgs (a -> b) = Succ (CountArgs b)
```

```
type instance CountArgs a          = Zero
```

Closed type families

Lets us write closed “type functions”:

```
type family CountArgs (f :: *) :: Nat where
  CountArgs (a -> b) = Succ (CountArgs b)
  CountArgs a         = Zero
```

Matching from top to bottom, like value-level functions.

Lets us write overlapping patterns, something we cannot do with open families.

Overloaded lists

```
class IsList l where
  type family Item l :: *
  fromList  :: [Item l] -> l
  toList    :: l -> [Item l]
  fromListN :: Int -> [Item l] -> l
```

```
instance IsList (Vector a) where
  type instance Item (Vector a) = a
  fromList  = V.fromList
  toList    = V.fromList
  fromListN = V.fromListN
```

Minimal type class instances

class Bifunctor p where

bimap :: (a -> b) -> (c -> d) -> p a c -> p b d

bimap f g = first f . second g

first :: (a -> b) -> p a c -> p b c

first f = bimap f id

second :: (b -> c) -> p a b -> p a c

second = bimap id

Minimal type class instances

This instance emits no warnings, but every method invocation will result in an infinite loop:

```
class Bifunctor (,) where
```

Minimal type class instances

class Bifunctor p where

```
{-# MINIMAL bimap | (first, second) #-}
```

```
bimap :: (a -> b) -> (c -> d) -> p a c -> p b d
```

```
bimap f g = first f . second g
```

```
first :: (a -> b) -> p a c -> p b c
```

```
first f = bimap f id
```

```
second :: (b -> c) -> p a b -> p a c
```

```
second = bimap id
```

Typed TemplateHaskell

- `TExp a`, like `Exp` but indexed by the type of the contained expression.
- To introduce `TExps` we use typed quasiquotes (`[| | ... | |]`) or typed splices (`$$x`).

```
compose :: Q (TExp (b -> c)) -> Q (TExp (a -> b))
        -> Q (TExp (a -> c))
compose f g = [| | $$f . $$g | |]
```

We can safely compile and run `TExps` at runtime.

Safe GeneralizedNewtypeDeriving

- `GeneralizedNewtypeDeriving + TypeFamilies = unsafeCoerce`
 - In other words, GHC Haskell < 7.7 is unsound.
 - The problem: GND treats the `newtype C = C Int` as “the same” as `Int`, but they might have different instances for a type family.
 - Fixed in 7.8 using “roles”.
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But wait, there's more!

- Parallel builds with `ghc -j`
- Improved IO manager
- Type natural solver
- `-XNumDecimals`
- Dynamic by default
- `clang` support
- iOS support and cross compilation

A more complete picture at <https://ghc.haskell.org/trac/ghc/wiki/Status/Oct13>

Future plans

- `class Applicative m => Monad m`
 - Pattern synonyms
 - Explicit type application
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